Computing period matrices of algebraic curves

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Definition

Let C be a smooth irreducible projective curve of genus g and let J be its Jacobian variety. Over the complex, J has the structure of a complex torus

$$J(\mathbb{C}) \cong \mathbb{C}^g / \Lambda,$$

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where Λ is a full rank lattice, called *period lattice* of C. For every basis $\omega_1, \ldots, \omega_g$ of the space of holomorphic differentials Ω^1_C we have that

$$\Lambda \cong \left\{ \int_{\gamma} \bar{\omega}, \gamma \in H_1(\mathcal{C}, \mathbb{Z}) \right\} \subset \mathbb{C}^g,$$

where $\bar{\omega} = (\omega_1, \dots, \omega_g)$ and $H_1(\mathcal{C}, \mathbb{Z}) \cong \mathbb{Z}^{2g}$ is the first homology group of the curve.

Definition (cont.)

Choosing a symplectic basis α_i , β_j $(1 \le i, j \le g)$ of $H_1(\mathcal{C}, \mathbb{Z})$, we define the matrices

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such that $\Lambda = \Omega \mathbb{Z}^{2g}$. We obtain a *small period matrix* in the Siegel upper half-space via

$$\tau = \Omega_A^{-1} \Omega_B \in \mathfrak{H}_g.$$

Period matrices or more generally, integration of differential forms on a Riemann surface, are required for computing

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• the Abel-Jacobi map

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- the endomorphism ring End(J) (numerical approximation),
- the real period of J (appearing in the BSD conjecture),
- the regulator pairing for K₂ of curves.

Existing work

For genus 1,2 and 3 period matrices can be computed in almost linear time to arbitrary precision (AGM, Borchardt mean).

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For hyperelliptic curves of arbitrary genus there exist

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 - a Maple implementation due to Deconinck and van Hoeij,
 - a Python/Sage implementation due to Swierczewski,
 - a Matlab implementation due to Frauendiener and Klein,
 - a Sage implementation due to Bruin is in progress.

Starting from an affine equation for the curve

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$$f(x,y) = 0,$$

we obtain a period matrix by working through the following list:

- computing a basis of holomorphic differentials
- choosing integration path \rightarrow analytic continuation
- numerical integration
- use the monodromy to compute a homology basis
- compute intersection matrix and symplectic base change

Our work (soon available)

Magma implementation for general algebraic curves (A1):

- based on the approach of Deconinck and van Hoeij,
- computes differentials using Magma's function fields,
- uses spanning tree methods to construct paths,
- analytic continuation is done via root approximation methods,

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Compared to the Maple implementation, we compute period matrices

- much faster and more reliably,
- to higher precision,
- for higher genera.

New algorithm for supereliptic curves

Consider a superelliptic curve given by an equation of the form

$$\mathcal{C}: y^m = f(x),$$

where $m \ge 2, \deg(f) \ge 3$ and $f \in \mathbb{C}[x]$ is separable.

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For such curves (in joint work with Pascal Molin) we developed and implemented algorithms in Magma and Arb (A2) that *rigorously* compute period matrices and the Abel-Jacobi map. More precisely:

- 'arbitrary' precision (realistically ≈ 10000 digits)
- excellent scaling with the genus (g >> 1000 possible)
- extremely fast and numerically robust
- better than Magma for hyperelliptic curves

Timings

Computation* of $\tau\in\mathfrak{H}_g$ for the family of curves given by

• $f_n = (x+y)^{n-1} + x^n y^2 + 1$ up to 20 significant digits

n	2	3	4	5	6	7	8	9	10
g	1	2	6	10	14	21	28	35	45
<i>t</i> _{Maple}	2.1s	бs	39s	2m 10s	error	6m 45s	12m 58s	-	error
t _{A1}	0.3s	0.8s	4.5s	15s	44s	2m 22s	5m 14s	12m 37s	30m 45s

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$$f_{m,n} = f_{m,n}$$

 $y^m - \sum_{k=0}^n x^k$ up to 500 significant digits

(m, n)	(2,5)	(2,11)	(2,31)	(2,101)	(3,5)	(3,11)	(7,5)	(77,5)	(11,21)	(31,21)
g	2	5	15	50	4	10	12	152	100	300
t _{A1}	32s	2m 30s	33m	-	1m 4s	5m 27s	4m 36s	-	38m	-
t _{A2}	0.2s	0.6s	3.7s	39s	4.8s	15s	6.7s	1m 26s	2m 4s	11m 14s
t _{Magma}	1.6s	6.7s	1m 23s	-	/	/	/	/	/	/

*done on Intel Xeon(R) CPU E3-1275 V2 3.50GHz processor.